

ECONOMICALLY SENSIBLE ALLOCATION OF RESOURCES FOR OVERLAY ROUTING RELAY NODES

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ABSTRACT:

The Internet is widely regarded as the most popular application on a worldwide scale. The major goal of the Internet is to provide complete services and a diverse range of information and communication capabilities, all while maintaining a high level of accessibility and responsiveness. Diverse and continuous efforts are being undertaken to identify and remedy issues affecting the accessibility of the Internet. Connection and router issues, as well as a variety of other issues, have the greatest impact on Internet performance. One potential option for improving Internet behavior is the widespread deployment of intelligent routers throughout network infrastructure, which would efficiently manage and control traffic flow by interacting with other intermediary devices. Intelligent nodes, also known as relay nodes or overlay nodes, allow the traversal of overlay paths, resulting in the formation of an overlay network. The aforementioned nodes have dynamic operational capabilities and add functionality to the existing infrastructure. Overlay Routing Resource Allocation (ORRA) is a versatile architecture that can be deployed to a wide range of applications while requiring less maintenance than other overlay systems such as Resilient Overlay Network (RON) and Detour.

Keywords: overlay, routing, relay

1. INTRODUCTION

The Internet has a large user base of millions of people, with thousands of new users joining every day. Millions of Internet users are served by interconnected networks with varied architectures. The networks are made up of autonomous systems that run independently, with distinct administrators in charge of maintaining and adjusting each one. The Internet's varied composition and proclivity for development highlight the need to improve its resilience, adaptability, efficiency, and accessibility. The Internet has several challenges, including software and hardware problems, performance concerns, congestion, and slow speed, to name a few. Many years have been spent investigating the inefficiencies of the Internet, and efforts have been made to improve its dependability and performance. There are numerous pathways

connecting the origin and destination. To select a single path for data transmission, multiple routing algorithms are used. It is crucial to note, however, that the route picked by the routing method may be longer in some cases. The overlay technique uses a selection mechanism to discover the shortest path from a group of multiple data paths, allowing for efficient and quick transmission. The overlay network employs a selection technique to locate and employ certain nodes inside the existing infrastructure that serve as data transmission intermediaries. It is critical to maintain and configure these intermediate nodes in order to maximize efficiency.

The overlay network boosts network performance and adds functionality to the existing network design without requiring any changes. Standardization and global deployment of new features may not be necessary. The inclusion of a virtual layer of functionality is notable. These intermediate nodes are known as relay nodes, and

they are specifically constructed for this purpose. Several overlay networks have been developed in the past, each delivering a unique set of services. These are application-specific and provide unique network capabilities. Overlay Transmission Control Protocol (oTCP) increases Transmission Control Protocol (TCP) performance by separating a TCP connection with a relatively long Round Trip Time (RTT) into many TCP connections with shorter RTT values. Using the Resilient Overlay Network (RON) improves the dependability of network systems. The RON network's numerous RON nodes are joined using a mesh topology, which is used to improve the network's fault tolerance. The tunneling concept is being used to boost resilience. To conceal specific data, the process entails encapsulating a packet within another packet and transferring it via a separate Internet channel. To accommodate client demand, the Global Internet Service Provider (G-ISP) employs a second Internet Service Provider (ISP).

To improve performance and functionality, carefully select overlay nodes from the existing infrastructure nodes. The implementation and maintenance of these overlay nodes entail major expenditures in addition to hardware, software, operations, and human engagement. As a result, it is critical to undertake a thorough study of the costs associated with improved performance and to choose a solution that is cost-effective. The costs of implementing previously investigated overlay networks were not considered. Multiple overlay networks with distinct services can coexist on the same substrate at the same time. All of these computations are about allocating resources. It is critical that the substrate distributes its resources in an equitable manner. Regardless of the application, the goal of ORRA is to determine the smallest number of infrastructure nodes required to maintain a specific overlay routing property. To achieve the shortest path routing strategy, the suggested method use BGP-based routing. Its goal is to find the smallest number of relay nodes required to construct the

shortest pathways between a group of autonomous systems. By employing a greedy technique to propose a limited number of relay nodes, BGP routing latency is reduced and overall performance is improved.

2. RELATED WORK

For more than three decades, TCP has been used to carry data across a wide range of applications. In the context of TCP, a decrease in Round-Trip Time (RTT) between connection endpoints results in a faster rise in window size, hence enhancing transmission. Pucha et al. (year) proposed Overlay TCP (oTCP) as an application-level protocol that functions as a TCP extension within the context of. This approach comprises separating a long round-trip time (RTT) TCP connection into several sub-connections with lesser RTTs. This is achieved by introducing intermediary nodes that increase throughput. oTCP is capable of detecting alternate channels that outperform the direct path in addition to reducing failures. For applications requiring high throughput, the oTCP architecture is a suitable choice. This is achieved by extending the existing TCP protocol without changing its behavior and terminating the overlay connection at the transport layer rather than the network layer. The letter E signifies a collection of connections that reflect communication relationships between entities. Let P_u denote the set of routing paths deduced from the underlying routing policy, and P_o denote the set of routing paths deduced from the overlay routing scheme. All direct routes with a single intermediate connection are found in the P_u region. The term P_o refers to the collection of the shortest pathways given by a weight function $W: E \rightarrow R$ over the vertices. This expression alludes to the cost of processing or the level of desire for a specific item's perfection. When a source-destination pair (s, t) is given, the overlay path between s and t , abbreviated $P_{s,t}$, is made up of two overlay pathways that connect s and t .

D. RON, a more stable application layer overlay network, was created by Andersen et al. RON nodes are drawn from several routing domains and

work together to ensure efficient data delivery. RON can detect and recover from path disturbances in less than twenty seconds. RON reduces data transmission loss, latency, and throughput by implementing alternate pathways in the event of failure detection and recovery. To create inter-node communication and discover network faults, the RON nodes use a proactive approach, employing aggressive probing and path monitoring techniques. The exchange of path quality information enhances data transmission by allowing the most optimal path for data transfer to be selected. For data transfer, the most efficient underlying Internet path is used. However, RON nodes are used to forward data if the underlying path is not ideal. To promote efficient and rapid system fault recovery, the RON nodes are interconnected in a comprehensive graph architecture. The costs associated with building RON are relatively expensive. Using data forwarding and failure recovery procedures, Redundant Overlay Networks (RON) improve application dependability and accessibility.

A routing system's principal function is to determine the best data transfer links. However, it has been noted that this strategy frequently results in an inconsistency in network traffic allocation, with certain routes remaining underutilized while others becoming overcrowded. Idle connections can be used to increase dependability and productivity. Savage et al. (year) developed the detour framework, which acts as a novel virtual network. This framework is designed to run on top of the existing Internet infrastructure [4]. Detour is built on a network of geographically separated router nodes linked together by a virtual connection known as a tunnel. Every packet entering a tunnel is encapsulated within a new IP packet before being sent across the Internet to the tunnel's predetermined exit point. Detour nodes are peripheral devices that allow path information to be shared. They use alternate pathways rather than dedicated connections.

R. proposes the concept of a Global Internet Service Provider (G-ISP). In their paper, Cohen et

al. suggest an Internet Service Provider (ISP) that offers transit services to its customers via an overlay network. This feature adds protocol support to the standard Border Gateway Protocol (BGP). The major goal of this design is to address and minimize difficulties associated with inter-domain routing protocols, such as sluggish convergence, insufficient quality of service (QoS), and restricted support for multicast communication. The system connects to the client through IP tunneling and provides inter-AS multicast capabilities.

3. PROPOSED SYSTEM

The suggested method's major goal is to create a flexible framework that may be applied to a range of overlay applications. Furthermore, it is critical to analyze the financial implications of this deployment and make every attempt to reduce the associated costs. The ORRA approach provides a framework that effectively meets stated criteria.

$G = (V, E)$ depicts the network topology of a graph, where V represents the topology nodes that imitate routing nodes and s and t represent the endpoints of path p for each $p \in P$, s, t . The selection of links is determined by examining the minimal link weights, which serve as indicators of effective bandwidth. Consider the Relay nodes that are used for overlay routing between sources and destinations as represented by the set U . Using underlay pathways makes it easier to route packets from one relay node to another. As proposed by Rami Cohen et al., U covers the interval (s, t) if there is a path p in the set $P(s, t)$ produced by concatenating one or more underlying paths and the endpoints of each of these underlying paths are in the set U s, t . In the event of a failure, the primary purpose is to decide which set U requires the fewest relay nodes for data transfer. Certain conditions must be met to ensure that the overlay nodes take the shortest path between each source-destination pair. The problem of Enhanced Overlay Routing Resource Allocation (eh-ORRA) is formalized as follows:

Definition:- The goal of selecting a subset of overlay nodes U from V in the context of a given graph $G=(V,E)$ and a set Q representing all source-destination pairings is to effectively cover Q while also serving as relay nodes that facilitate data transfer over the shortest possible paths.

System Architecture

The system architecture is represented in Figure 1. In modern networking, BGP is the dominant inter-domain routing protocol. Border Gateway Protocol (BGP) allows BGP systems to share network reachability information. The shortest transmission paths can be chosen based on the path metrics. Certain network parts may be preserved in order to increase overall performance. These entities are commonly referred to as relay nodes or overlay nodes. These devices act as routers, allowing data packets to be sent to their intended destinations. When the sender commences data transmission, the overlay routers choose the best path depending on a variety of factors. Data transfer happens between the source and the overlay router, as well as between the overlay router and the destination. End-to-end semantics are preserved, and data is delivered to the intended recipient.

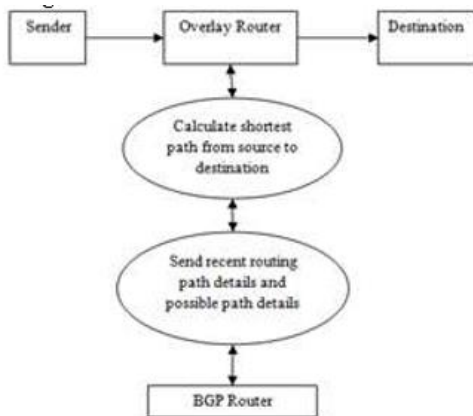


Figure1: System Architecture

Algorithm

The Overlay Routing Resource Allocation (ORRA) approach is used to identify potential relay nodes. The strategy used can be described as avaricious. The end result is a solution that is close to ideal. The following is a recursive algorithm for solving the NP-hard ORRA issue.

Terms:

$Q = \{(s_1, t_1), (s_2, t_2), \dots, (s_n, t_n)\}$ - all source destination pairs

P^s - set of overlay paths between s and t

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Inputs: $G(V,E)$ network to pology, W - Weight function $W: V \times V \rightarrow R^+$, P_U - Set of underlay paths, P_o - set of overlay paths, U - Set of relay nodes.

Steps:

ORRA($G=(V,E), W, P_u, P_o, U$)

Step 1: $\forall v \in (V-U)$, if $w(v)=0$ then $U \cup \{v\}$ Step 2: If U covers Q then return U

Step 3: Find $(s,t) \in Q$ not covered by U

Step 4: Find $\forall v$ present on P^s and $v \in U$ then $V \cup \{v\}$ Step 5: set $x = \{\min w(v) \mid v \in V\}$

Step 6: set $w_1(v) = \begin{cases} x, & v \in V \\ 0, & \text{otherwise} \end{cases}$

$0, \text{ otherwise}$

Step 7: $\forall v \in V$ set $w_2(v) = w(v) - w_1(v)$

Step 8: ORRA(G, w_2, P_u, P_o, U) goto step 1 Step 9: $\forall v$ if $U - \{v\}$ covers Q then $U = U - \{v\}$ Step 10: return U

Output: These to relay nodes U is the output of algorithm.

Functional Diagram

Figure 2 displays the many functions performed by the ORRA algorithm during relay node selection. In order to function properly, the method requires access to particular network topology information, which is provided as input. When all nodes understand the same network information, TCP connections are established between them. BGP is in charge of developing the routing policies that govern path selection. Multiple Border Gateway Protocol (BGP) routers exchange path information. Individual paths can be assigned different ratings by the Border Gateway Protocol (BGP) based on a range of parameters such as delay, bandwidth, and congestion. When network nodes need to send data, they must first decide on a source-destination pair. They choose between a direct path and an overlay route based on the path rating decision. In the event that an overlay path is chosen, relay nodes along that path are chosen to ensure data delivery. The procedure comprises identifying overlay nodes, which are then used for data transmission.

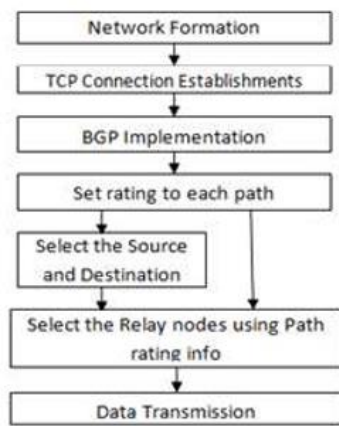


Figure2: The current investigation is focused on the functional diagram of the ORRA algorithm.

4.PERFORMANCE ANALYSIS

This paper presents a novel strategy for optimizing resource allocation in overlay networks by utilizing link and node weights. The proposed Link-Node weight based ORRA-L N approach is compared to the Node weight based ORRA-N technique in the evaluation. The ORRA-N approach allocates transmission resources exclusively depending on each node's weight. Unlike the traditional ORRA-LN technique, node weight and link weight are used to properly distribute resources for data transmission. The two key performance measures examined during the routing cost evaluation are throughput and latency.

Performance Metrics for Delay:

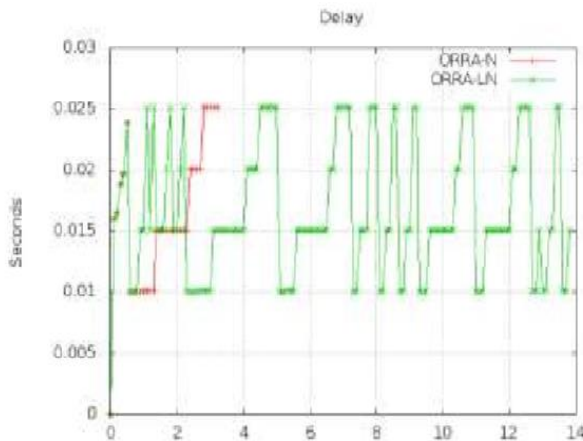


Figure3 The goal of this inquiry is to undertake a delay performance analysis. For the purpose of a test case, a simulation was

carried out utilizing both the ORRA-N and ORRA-NL techniques. The duration of the delay has been reduced slightly. The delay graph demonstrates the substantial variations in response times that have an influence on the proposed system. This occurrence occurs as a result of transit congestion or timeouts, causing the provider's intended route to execute poorly. Observing throughput gives factual evidence of measurable improvement.

Performance Metrics for Throughput:

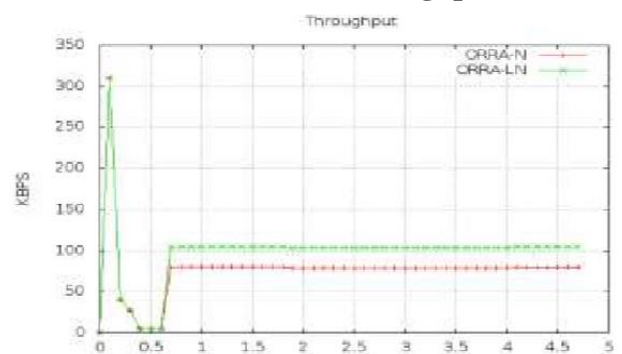


Figure4: Analysis of data transfer rate performance indicators, specifically throughput.

The graph of throughput shows a strong surge followed by a gradual drop. When a certain threshold is reached, the value remains constant. As throughput grows, the ORRA-LN technique outperforms the ORRA-N technology in terms of efficacy. However, as network complexity increases, both latency and throughput suffer as more resources are added to the overlay network. This eventually leads to a drop in average throughput. The proposed technology significantly improves throughput, resulting in cheaper transmission costs. This means it has a high level of efficacy at a low cost.

5. CONCLUSION

Previous overlay systems were limited to specific applications and did not account for the costs associated with the deployment of overlay nodes. The proposed Optimal Resource Allocation (ORRA) technique considers and strives to reduce deployment costs. Furthermore, it provides a complete foundation that may be applied to a wide

range of applications. The system provides a collection of relay nodes that can determine the shortest path while taking the weights assigned to each node into account. When link weights are taken into account, the resulting path is more cost-effective than when only node weights are taken into account.

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